

# CubeHash round-2 modifications

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This document describes changes from the round-1 CubeHash submission package to the round-2 CubeHash submission package.

**CubeHash specification (2.B.1)** ([spec.pdf](#)) defines  $\text{CubeHash}_{r/b-h}$  using exactly the same text as the original CubeHash submission;  $\text{CubeHash}_{r/b-h}$  is exactly the same function in round 2 that it was in round 1. However, the recommendations for parameters  $(r, b)$  have been updated as described in my note “CubeHash parameter tweak: 16 times faster”:

- $\text{CubeHash}_{16/32-224}$  is proposed for  $\text{SHA-3-224}$ ,
- $\text{CubeHash}_{16/32-256}$  is proposed for  $\text{SHA-3-256}$ ,
- $\text{CubeHash}_{16/32-384}$  is proposed for  $\text{SHA-3-384-normal}$ ,
- $\text{CubeHash}_{16/32-512}$  is proposed for  $\text{SHA-3-512-normal}$ ,
- $\text{CubeHash}_{16/1-384}$  is proposed for  $\text{SHA-3-384-formal}$ , and
- $\text{CubeHash}_{16/1-512}$  is proposed for  $\text{SHA-3-512-formal}$ .

There is also a new subsection “Additional comments on symmetries” extending the symmetry paragraph in the original submission.

**CubeHash efficiency estimates (2.B.2)** ([estimates.pdf](#)) now summarizes eBASH Core 2 Duo benchmarks for CubeHash, confirming the original efficiency estimates. The document also adds a paragraph discussing microarchitectural variability among Core 2 Duo CPUs and recommending that NIST specify which CPU is actually the reference platform.

**CubeHash expected strength (2.B.4)** ([strength.pdf](#)) has been modified to note the expected impact of quantum computers. Grover’s algorithm will find (e.g.) 224-bit preimages for any of the SHA-3 candidates in only about  $2^{112}$  quantum operations. This quantum computer

- has a much higher success chance than a conventional computer performing  $2^{200}$  operations and
- is much more likely to be available to future attackers than a conventional computer performing  $2^{200}$  operations,

so considering the conventional threat while ignoring the quantum threat makes no sense from a risk-analysis perspective.

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**CubeHash attack analysis (2.B.5)** ([attacks.pdf](#)) has been reorganized and expanded. The document includes the description of narrow-pipe attacks that had appeared in the original submission as “CubeHash appendix: complexity of generic attacks.” The document also reviews various third-party analyses of CubeHash that have been announced by Aumasson, Bloom, Brier, Dai, Janis, Kaminsky, Khazaei, Khovratovich, Meier, Naya-Plasencia, Nikolic, Peyrin, Rao, Salaev, Wang, Weinmann, and Wilson. The most recent third-party analysis is the Brier–Khazaei–Meier–Peyrin paper “Linearization framework for collision attacks: application to CubeHash and MD6” to appear at Asiacrypt 2009.

**CubeHash features (2.B.6)** ([features.pdf](#)) has been extended to include subsections “Unified implementation across output sizes,” “Small code size and vector-code size,” and “Good security/speed tradeoff.”